

Some complements on Kan complexes
(based on Ex. sheet 2).

Lemma 32: (Exercise 2.4)

X . Kan complex. The relation on X_0

given by: $x \sim y \Leftrightarrow \exists f \in X_1, \begin{cases} d_1(f) = x \\ d_0(f) = y \end{cases}$

is an equivalence relation.

$$\Rightarrow \pi_0(X) = X_0 / \sim.$$

proof: Reflexivity $d_0 s_0(x) = d_1 s_0(x) = x$ (does not require X . Kan)

Symmetry Assume $x \sim y$ witnessed by $p \in X_1$.

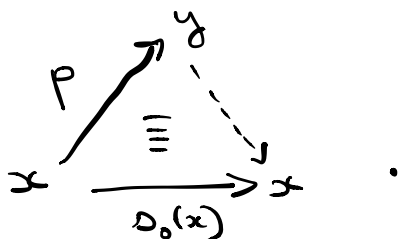
$$\text{We have } \Lambda_0^2 = \Delta^1 \coprod_{\Delta^0} \Delta^1, \text{ i.e. } \begin{array}{ccc} \Delta^0 & \xrightarrow{\langle 0 \rangle} & \Delta^1 \\ \langle 0 \rangle \downarrow & & \downarrow \\ \Delta^1 & \xrightarrow{\Gamma} & \Lambda_0^2 \end{array}$$

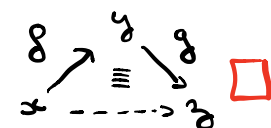
$$\text{so } \text{Set}(\Lambda_0^2, X) = \left\{ (p_1, p_2) \in X_1^2 \mid d_1(p_1) = d_1(p_2) \right\}$$

Consider $\Lambda_0^2 \rightarrow X$. given by $(p, s_0(x))$.

Apply the Kan lifting property $\Rightarrow \Delta^2 \xrightarrow{t} X$.

Then $d_0 t \in X_n$ witnesses $y \sim x$:



Transitivity Same construction with Δ^2_1 :  □

def 33: A simplicial set X_* is **discrete** if for all $g \in \Delta([m], [n])$, $g^x: X_n \rightarrow X_m$ is a bijection.

Lemma 33: (Exercises 1.6, 2.5)

a) The functors

$$\text{Set} \longrightarrow \text{sSet}^{\text{discr}}$$

$$S \longmapsto \underline{S} := \text{constant presheaf with value } S:$$

$$[n] \mapsto S, \quad g \mapsto \text{id}_S.$$

and $\text{sSet}^{\text{discr}} \xrightarrow{\text{ev}_{[0]}} \text{Set}$ yield an equivalence of categories between discrete simplicial sets and sets.

b) discrete simplicial sets are Kan complexes.

proof: a) We have $ev_0(\underline{\Delta}) \cong S$ naturally in S .

It remains to show $ev_0(X_\bullet) \cong X_\bullet$ for X_\bullet discrete.

There is actually a natural transformation $ev_0(-) \xrightarrow{\varepsilon} id$ of functors $sSet \mathcal{C}$, given on X_\bullet by:

$$(\varepsilon_{X_\bullet})_n: \underbrace{ev_0(X_\bullet)}_n = X_0 \xrightarrow{\langle 0 \dots 0 \rangle^*} X_n. \quad \left(\begin{array}{l} \text{This is a map of s-sets} \\ \text{because } [0] \text{ is the final obj. of } \Delta \end{array} \right)$$

If X_\bullet is discrete, then by definition ε_{X_\bullet} is an isomorphism.

and we are done. $(\underline{Rmk} \text{ Psh}(\mathcal{C})^{\text{discr}} \cong \text{Sets} + \text{action of } \pi_1(\text{INCL}))$

b) By part a) we can assume $X_\bullet = \underline{\Delta}$ (we could also argue directly)

We must show: $\forall n \geq 1, \forall 0 \leq k \leq n, sSet(\Delta^n, \underline{\Delta}) \cong S$

$$\downarrow \\ sSet(\Lambda_R^n, \underline{\Delta})$$

I will explain in general how to compute $sSet(X_\bullet, \underline{\Delta})$.

Claim 1: A map $X_\bullet \xrightarrow{f} \underline{\Delta}$ is determined by $X_0 \xrightarrow{f_0} S$.

Pf: $\xrightarrow[n]{\langle 0 \dots 0 \rangle^*} f_n(x) = f_0(\xrightarrow[n]{\langle 0 \dots 0 \rangle^*} x)$; but $\langle 0 \dots 0 \rangle^*$ is a bijection in $\underline{\Delta}$.

Claim 2: $f_0: X_0 \rightarrow S$ factors through $X_0 \xrightarrow{\quad} S$
 $\searrow \pi_0(X_\bullet) \dashrightarrow$

Pf: Let $x, y \in X_0$ and $p \in X_1$, $d_1(p) = x$ and $d_0(p) = y$.

$$\text{Then } \begin{cases} d_1 \circ f_1(p) = f_0(d_1 p) = f_0(x) \\ d_0 \circ f_1(p) = f_0(d_0 p) = f_0(y) \end{cases} \text{ but in } \underline{S}, d_0 = d_1 = \text{id}_S !$$

$$\Rightarrow f_0(x) = f_0(y)$$

This is enough to finish the exercise; indeed $\pi_0(\Lambda_R^n)$ has one element (because $I^n \subset \Lambda_R^n$, hence $0 \cup 1 \cup \dots \cup n$). ($n \geq 2$)

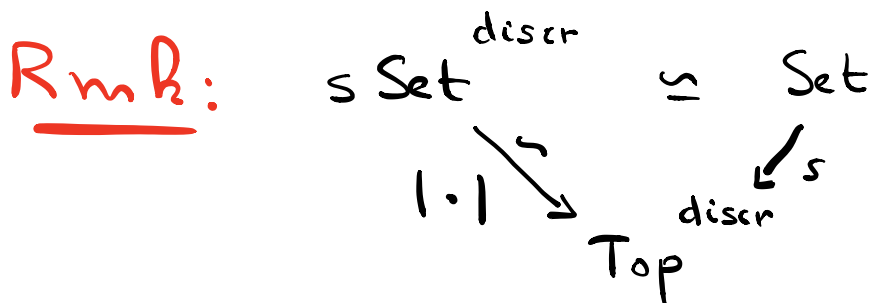
Claim 3: $s\text{Set}(X, \underline{S}) \cong \text{Set}(\pi_0(X), S)$.

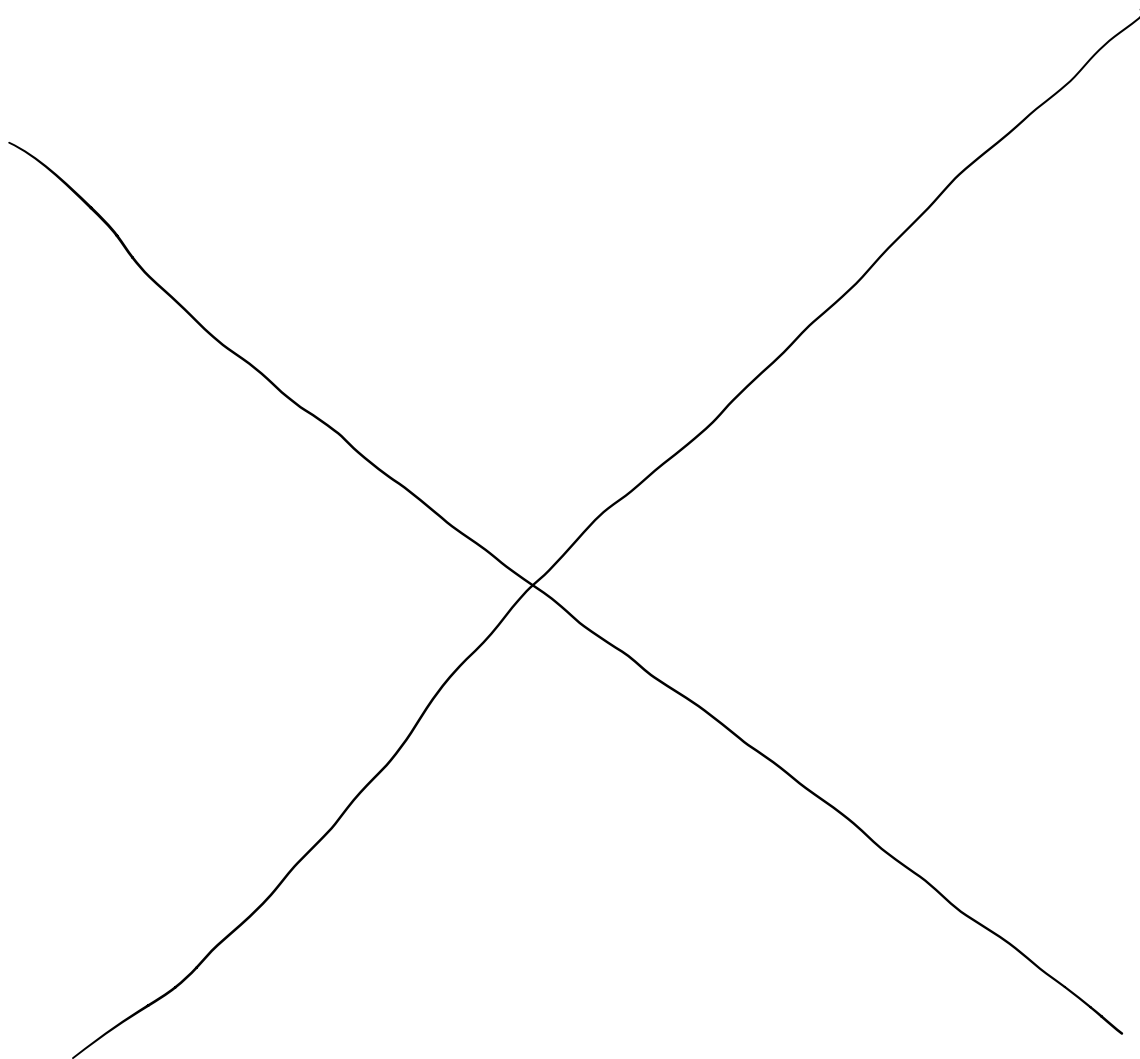
Pf: We have constructed an injective map \hookrightarrow .

It remains to see that any map $\pi_0(X) \xrightarrow{g} S$ can be realized this way. There is a natural candidate:

$$x \in X_n \longmapsto g([\langle 0 \dots 0 \rangle^* x]) \in S$$

I leave it to you to check that this is well-defined and gives an inverse. □





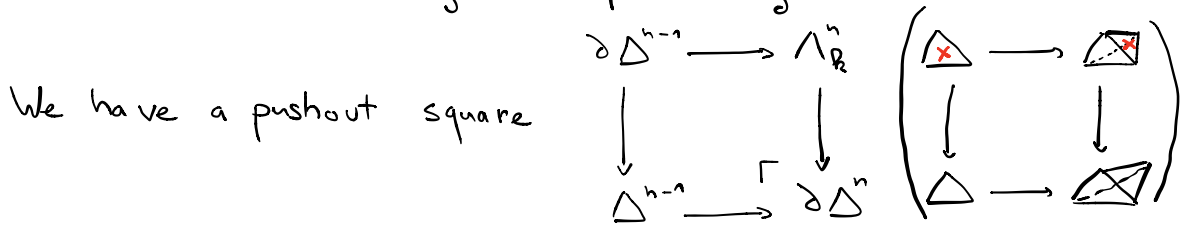
Lemma 34: (Exercise 2.6)

Trivial Kan fibrations are Kan fibrations.

proof: Let $f: X \rightarrow Y$ be a trivial Kan fibration, and $0 \leq R \leq n$. We are given a diagram:

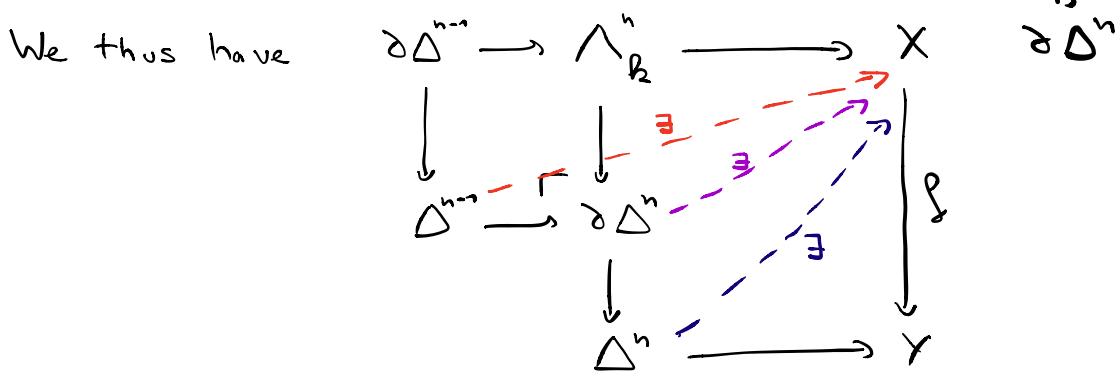
$$\begin{array}{ccc} \Lambda_R^n & \rightarrow & X \\ \downarrow & & \downarrow \\ \Delta^n & \rightarrow & Y \end{array}$$


and look for a diagonal map filling it.





To prove it is a pushout, can use Prop 24 on

the skeletal filtration $S\mathbb{R}_{n-2}(\partial \Delta^n) \hookrightarrow S\mathbb{R}_{n-1}(\partial \Delta^n)$



with  exists because f is a trivial Kan fibration

 exists because of the universal property of the pushout square.

 exists because f is a trivial Kan fibration.



II) Infinity - categories

1) Nerves of categories

Our model of ∞ -categories will be a particular type of simplicial sets. We also want to be able to consider ordinary categories as ∞ -categories.

\rightsquigarrow We need a fully faithful functor

$$N: \text{Cat} \hookrightarrow \text{sSet} .$$

- Posets give rise to categories:

$$\text{Poset} \xrightarrow{\text{fully faithful}} \text{Cat}$$

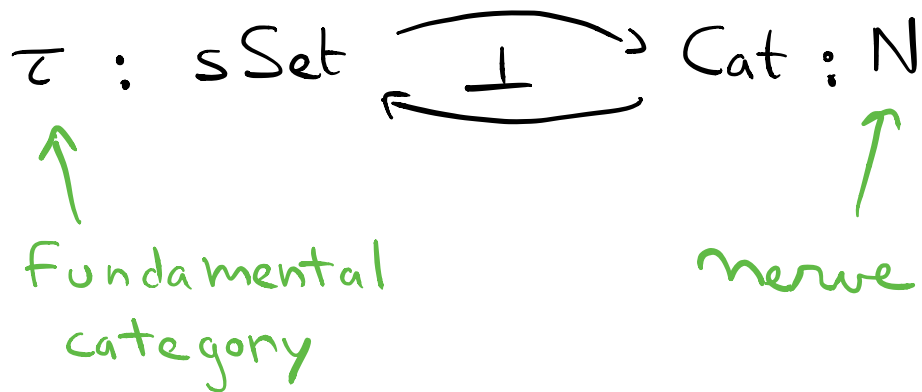
$$P \longmapsto \text{Ob} : x \in P$$

$$P(x, y) = \begin{cases} *, & x \leq y \\ \emptyset, & \text{otherwise} \end{cases}$$

In particular we have a fully faithful cosimplicial category:

$$\Delta \xrightarrow{Q^\bullet} \text{Cat}$$

def 1 By the free cocompletion property, Q^\bullet induces an adjunction:



By construct^o:

$$\left\{ \begin{array}{l} \tau(\Delta^n) = [n]. \\ \tau \text{ commutes with colimits.} \end{array} \right.$$

and

$$N(C)_n = \text{Cat}([n], C) \left| \begin{array}{l} \text{sSet}(\Delta^n, N(C)) \\ \text{is} \\ \text{Cat}([n], C) \end{array} \right.$$

$$= \left\{ \begin{array}{l} n\text{-tuples of} \\ \text{composable morphisms} \end{array} \right\}$$

eg: $(NC)_0 = \text{Ob } C$, $(NC)_1 = \text{Mor } C$

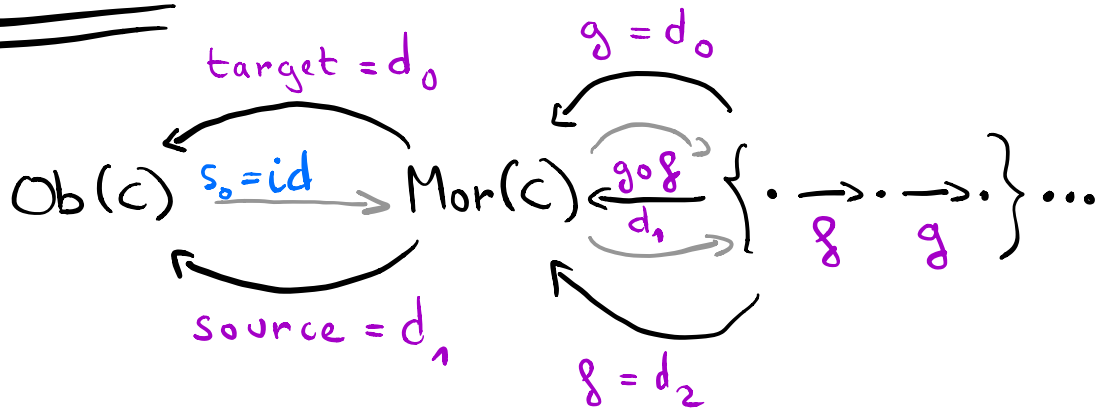
$$(NC)_n = \left\{ \cdot \xrightarrow{\delta_1} \cdot \xrightarrow{\delta_2} \dots \xrightarrow{\delta_n} \cdot \text{ in } C \right\}$$

Rmk We have $N([n]) \simeq \Delta^n$

($\Leftarrow Q$ fully faithful).

$$N([n])_m = \text{Cat}([m], [n]) \stackrel{\downarrow}{=} \Delta([m], [n])$$

$N(C)$:



This makes it clear that

we can reconstruct C from $N(C)$.

Prop 2 N is fully faithful

proof: For C, D in Cat , we must show

$$\text{Cat}(C, D) \xrightarrow{\sim} \text{sSet}(NC, ND).$$

Injectivity: A functor is determined by its effect on objects and morphisms. Since $(NC)_0 = \text{Ob}(C)$ and $(NC)_1 = \text{Mor}(C)$, we are done.

Surjectivity Let $\alpha: NC \rightarrow ND$.

We define a candidate F for the preimage functor using again

$$\left\{ \begin{array}{l} (NC)_0 = \text{Ob}(C) \\ (NC)_1 = \text{Mor}(C) \end{array} \right. \left| \begin{array}{l} F|_{\text{Ob}} = \alpha_0 \\ F|_{\text{Mor}} = \alpha_1 \end{array} \right.$$

To check that F is a functor, we need to see:

- F compatible with source/target:

Let $f: c \rightarrow c'$ in C

$$\begin{cases} s F(f) \stackrel{\text{def}}{=} d_1 \alpha(f) \stackrel{\text{sSet}}{=} \alpha d_1(f) \stackrel{\text{def}}{=} F(c) \\ t F(f) \stackrel{\text{def}}{=} d_0 \alpha(f) \stackrel{\text{sSet}}{=} \alpha d_0(f) \stackrel{\text{def}}{=} F(c') \end{cases}$$

- F compatible with identities

$$F(\text{id}_c) \stackrel{\text{def}}{=} g s_0(c) \stackrel{\text{sSet}}{=} s_0 g(c) \stackrel{\text{def}}{=} \text{id}_{F(c)}$$

- F compatible with compositions

$$\begin{aligned} F(g \circ f) &= \alpha(g \circ f) = \alpha(d_1(g, f)) \\ &= d_1 \alpha(g, f) = F(g) \circ F(f) \end{aligned}$$

$\leadsto F$ is a functor.

By construction, $N(F)$ and α coincide on 0 and 1-simplices, and it is easy to see that it forces them to be equal

(because simplices in $(NC)_i$ for $i > 2$ are uniquely determined by their 1-simplices). \square

Example M monoid \leadsto 1-object category BM

$N(BM)$ classifying simplicial set: $N(BM)_n = M^{x_n}$
 $|N(BM)|$ ————— space of $M \leadsto$ group cohomology

- We also need a precise description of fundamental categories.

Prop 3: Let $X_0 \in \mathcal{S} \text{Set}$. The fundamental category τX_0 admits a presentation by generators and relations:

- $\text{Ob } \tau X_0 = X_0$.
- $\text{Mor } \tau X_0$ is generated by X_n ; for all $n \geq 0$ and $f_1, \dots, f_n \in X_n$ such that $d_1 f_i = d_0 f_{i-1}$, we have $\overline{f_n} \circ \dots \circ \overline{f_1} \in \text{Mor } \tau X_0 (d_1 f_n, d_0 f_1)$.
- We have relations:
 - $\forall x \in X_0, \overline{\Delta_0(x)} = \text{id}_x \in \tau X_0(x, x)$;
 - $\forall t \in X_2^{(\text{nd})}$, we have $\overline{d_0 t} \circ \overline{d_2 t} = \overline{d_1 t}$.

In other words, for any other category D , we have $\text{Cat}(\tau X_0, D)$ is in natural bijection with the set of pairs

$$\left(\begin{array}{l} F_0: X_0 \longrightarrow \text{Ob } D, \\ F_1: X_1 \longrightarrow \text{Mor } D \end{array} \right) \quad \text{such that:}$$

$$\cdot \forall f \in X_1, \begin{cases} \text{source}(F_1(f)) = F_0(d_1 f) \\ \text{target}(F_1(f)) = F_0(d_0 f) \end{cases}$$

$$\forall x \in X_0, F_1(d_0 x) = \text{id}_{F_0(x)}$$

$$\cdot \forall t \in X_2^{(nd)}, F_1(d_1 t) = F_1(d_0 t) \circ F_1(d_2 t)$$

proof: We use the skeletal filtration together with the facts that $\begin{cases} \tau \text{ commutes with colimits.} \\ \tau(\Delta^n) = [n]. \end{cases}$

$$\cdot \tau(X_\bullet) = \text{colim}_{n \geq 0} \tau(\text{sk}_n X_\bullet)$$

category with one morphism.

$$\cdot \tau(\text{sk}_0 X_\bullet) = \tau\left(\coprod_{x \in X_0} \Delta^0\right) = \coprod_{x \in X_0} [0] = \coprod_{x \in X_0} *$$

$$\cdot \partial \Delta^1 = \Delta^0 \amalg \Delta^0 \Rightarrow \tau \Delta^1 = * \amalg *$$

$$\begin{array}{ccc} \coprod_{X_1^{nd}} \partial \Delta^1 & \longrightarrow & \coprod_{X_1^{nd}} \Delta^1 & \xrightarrow{\quad} & \coprod_{X_1^{nd}} (* \amalg *) & \longrightarrow & \coprod_{X_1^{nd}} [1] \\ \downarrow & & \downarrow & \Rightarrow & \downarrow & & \downarrow \\ \text{sk}_0(X_\bullet) & \longrightarrow & \text{sk}_1(X_\bullet) & & \coprod_{X_0} * & \longrightarrow & \tau(\text{sk}_1(X_\bullet)) \end{array}$$

$\Rightarrow \tau(\text{sk}_1(X_\bullet))$ is the free category generated

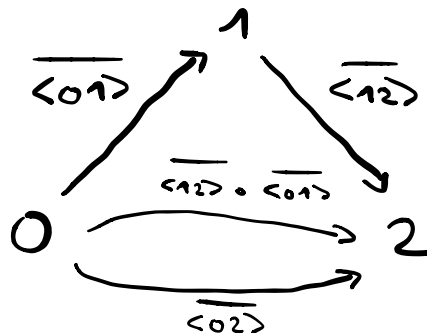
by X_1^{nd} .

Since $X_1 = X_1^{nd} \amalg_{S_0(X_0)}$, this is equivalent to generating by X_1 and imposing $\overline{s_0(x)} = id_x$.

• $\partial \Delta^2 = Sh_1(\partial \Delta^2)$

\Downarrow

$\tau(\partial \Delta^2)$ is given by:



• So the map $\tau(\partial \Delta^2) \rightarrow \tau(\Delta^2) = [2]$ identifies the two morphisms between 0 and 2.

• $\amalg_{X_2^{nd}} R(\partial \Delta^2) \longrightarrow \amalg_{X_2^{nd}} [2]$

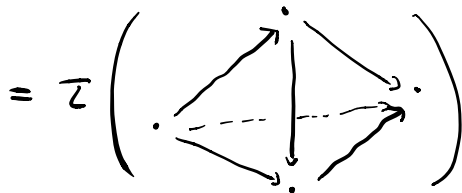
thus imposes exactly the relations of the statement for $t \in X_2^{nd}$.

$\downarrow \qquad \quad \quad \quad \downarrow$
 $\tau(Sh_1 X) \longrightarrow \tau(Sh_2 X)$

• Any $t \in X_2^{deg}$ imposes a relation with identities which is already there.

• $\tau(\partial \Delta^3) = \tau(Sh_2 \partial \Delta^3)$

($\partial \Delta^3$ is 2-skeletal)



= [3]

(use previous step)

$$\Rightarrow \tau(\partial\Delta^3) \xrightarrow{\sim} \tau(\Delta^3)$$

$$\Rightarrow \forall X \in \mathbf{sSet}, \tau(Sk_2 X) \xrightarrow{\sim} \tau(Sk_3 X)$$

skeletal
filtration

$$\Rightarrow \tau(\partial\Delta^4) \cong \tau(Sk_2 \partial\Delta^4) \cong \tau(\Delta^4)$$

$$\Rightarrow \forall X \in \mathbf{sSet}, \tau(Sk_3 X) \xrightarrow{\sim} \tau(Sk_4 X)$$

⋮

$$\Rightarrow \forall X \in \mathbf{sSet}, \tau(X) \xrightarrow{\sim} \tau(Sk_2 X)$$



Rmk The terminology "fundamental category" comes from topology:

- For $A \in \mathbf{Top}$, $\tau(\text{Sing } A) \cong \pi_{\leq 1} A$ fundamental groupoid of A .

- For $X \in \mathbf{Kan}$, $\tau(X) \cong \pi_{\leq 1} |X| \xrightarrow{\quad} |X|$

Notation $X. \in \mathbf{sSet}$, $x \in X_n$, $\langle g_0 \dots g_m \rangle : [m] \rightarrow [n]$.

We write $x_{g_0 g_1 \dots g_m} := \langle g_0 \dots g_m \rangle^*(x) \in X_m$

In particular, x_0, \dots, x_n are the vertices of x
 $x_{i,j}$ are the edges, and so on.

Thm 3 : (characterisation of nerves)

Let $X_\bullet \in \mathbf{sSet}$. TFAE:

(1) $X_\bullet \cong N\mathcal{C}$ for some $\mathcal{C} \in \mathbf{Cat}$.

(2) The unit $X_\bullet \xrightarrow{\eta} N \tau X_\bullet$ is an isomorphism.

(3) X_\bullet satisfies the unique

Spine extension property:

For all $n \geq 2$,

$$I^n \longrightarrow X_\bullet$$

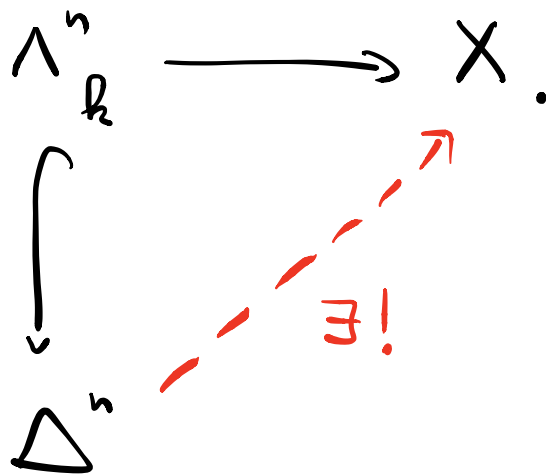
$$\downarrow$$
$$\Delta^n$$

$$\dashrightarrow \exists!$$

Grothendieck-Segal property.

(4) X satisfies the inner
horn unique extension property:

For all $n \geq 2$ and $0 < k < n$:



proof: (2) \Rightarrow (1) \checkmark

(1) \Rightarrow (2): Assume $X \cong NC$.

By a general property of
adjunctions, we have

$$\begin{array}{ccc}
 NC & \xrightarrow{\eta^N} & N \tau NC \\
 & \searrow & \downarrow N\varepsilon \\
 & & NC
 \end{array}
 \quad (*)$$

• Because N is fully faithful,

$\varepsilon : \tau NC \xrightarrow{\sim} C$ is an iso

$\Rightarrow N\varepsilon$ iso $\stackrel{(*)}{\Rightarrow} \eta^N$ iso

$\Rightarrow X \xrightarrow[\eta]{\sim} N \tau X$.

(1) \Rightarrow (3):

$\text{SSet}(\Delta^n, N(C))$

$$\left\{ c_0 \xrightarrow[\delta_1]{\text{IS}} c_1 \rightarrow \dots \rightarrow c_n \text{ in } C \right\}$$

IS

$$\left\{ f_1, \dots, f_n \in N(C)_1 \mid d_0(f_i) = d_1(f_{i+1}) \right\}$$

$$IS \leftarrow I^n = \Delta^{\{0,1\}} \underset{\Delta^{\{1\}}}{\parallel} \dots \underset{\Delta^{\{n-1\}}}{\parallel} \Delta^{\{n-1,n\}}$$

$$S\text{Set}(I^n, N(C))$$

(3) \Rightarrow (1):


We have to construct a category C and an isomorphism $X \xrightarrow{\sim} NC$.

We put $\text{Ob } C = X_0$, $\text{Mor } C = X_1$.

with $\text{source}(f) = d_1(f)$, $\text{target}(f) = d_0(f)$.

We put $\text{id}_x := s_0(x)$.

Let $f, g \in X_1$, $f: x \rightarrow y$ and $g: y \rightarrow z$.

we get $(f, g): I^2 \rightarrow X$. 

By assumption there is a unique extension t

to Δ^2 , and we define $g \circ f$ to be $d_1(t)$:

$$\begin{array}{ccc} & y & \\ f \nearrow & & \searrow g \\ x & t & z \\ & \text{---} & \\ & g \circ f & \end{array}$$

The unitality then follows from the triangles

$$s_0(g) : \begin{array}{ccc} & x & \\ s_0(x) \swarrow & & \searrow g \\ x & \xrightarrow{g} & y \end{array} \quad \text{and} \quad s_1(g) : \begin{array}{ccc} & y & \\ g \nearrow & & \searrow s_0(y) \\ x & \xrightarrow{g} & y \end{array}$$

Associativity follows from the uniqueness of liftings along $I^3 \hookrightarrow \Delta^3$.

\Rightarrow we have constructed a category C .

• We construct a morphism $X \xrightarrow{\psi} NC$.

Let $x \in X_n = s\text{Set}(\Delta^n, X)$. Then

$x|_{I^n}$ determines a sequence of composable morphisms in $C \rightarrow$ a simplex $\psi(x) \in (NC)_n$.

It is easy to see that $X \xrightarrow{\psi} NC$ is indeed a morphism of simplicial sets,

$$\text{with } \begin{cases} X_0 \xrightarrow{\cong} (NC)_0 \\ X_n \xrightarrow{\cong} (NC)_n \end{cases}$$

$$\begin{array}{ccc} s\text{Set}(\Delta^n, X) & \xrightarrow{\psi_n} & s\text{Set}(\Delta^n, NC) \\ \downarrow & & \downarrow \\ s\text{Set}(I^n, X) & \longrightarrow & s\text{Set}(I^n, NC) \end{array}$$

We now contemplate the diagram:

$$\begin{array}{ccc} X_n & \longrightarrow & (NC)_n \\ \downarrow S(3) & & \downarrow S((1) \Rightarrow (3)) \end{array}$$

$$X_n \times_{X_0} X_n \times \dots \times X_n \xrightarrow{\sim} (NC)_0 \times \dots \times (NC)_1$$

and see that $X_n \xrightarrow{\sim} (NC)_n$

(1) \Rightarrow (4):

Let $n \geq 2$ and $0 < k < n$.

We compute:

$$\begin{aligned} & \text{SSet}(\Lambda_k^n, N(C)) \\ & \simeq \text{Cat}(\tau(\Lambda_k^n), C) \end{aligned}$$

$$\cong \text{Cat}(\tau(\text{Sk}_2(\Lambda_R^n)), \mathbb{C})$$

We now observe that

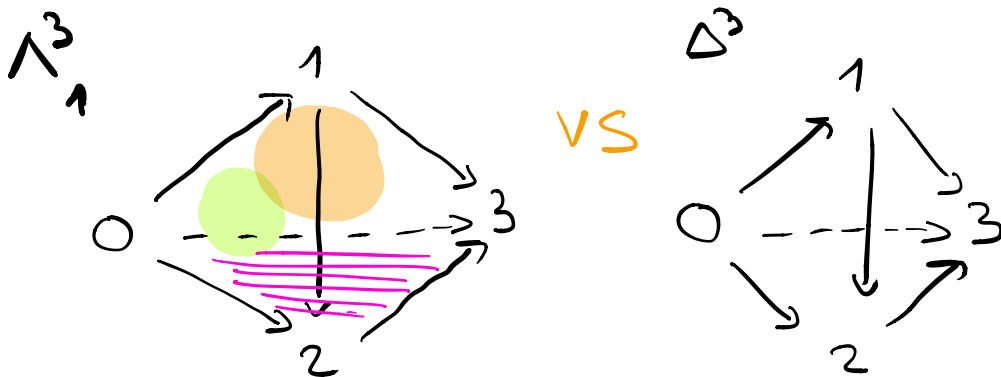
$$\begin{cases} \forall n \geq 4, \text{Sk}_2(\Lambda_R^n) \xrightarrow{\sim} \text{Sk}_2(\Delta^n) \\ \Lambda_1^2 = \mathbb{I}^2 + (1) \Rightarrow (3). \end{cases}$$

So it remains to show

$$\begin{cases} \tau(\Lambda_1^3) \xrightarrow{\sim} \tau(\Delta^3) \\ \tau(\Lambda_2^3) \xrightarrow{\sim} \tau(\Delta^3) \end{cases}$$

The two cases are similar, we

do Λ_1^3 . First a picture:



Λ^3_1 and Δ^3 have the same objects and edges. The only additional relation in $\tau(\Delta^3)$ is

$$\overline{\langle 03 \rangle} = \overline{\langle 23 \rangle} \circ \overline{\langle 02 \rangle}$$

but in $\tau(\Lambda^3_1)$ we have

$$\overline{\langle 03 \rangle} = \overline{\langle 13 \rangle} \circ \overline{\langle 01 \rangle}$$

back
face



$$= (\overline{\langle 23 \rangle} \circ \overline{\langle 12 \rangle}) \circ \overline{\langle 01 \rangle}$$

$$= \overline{\langle 23 \rangle} \circ (\overline{\langle 12 \rangle} \circ \overline{\langle 01 \rangle})$$

left
front
face



$$= \overline{\langle 23 \rangle} \circ \overline{\langle 02 \rangle}$$

$$\triangle \begin{cases} \tau(\Lambda_0^3) \neq \tau(\Delta^3) \\ \tau(\Lambda_3^3) \neq \tau(\Delta^3) \end{cases}$$

$\Rightarrow N(C)$ does not have a lifting property for outer horns.

(4) \Rightarrow (1):

We follow the same strategy as for (3) \Rightarrow (1). We construct a category C in exactly the same way, except that associativity is proved using

$$\Lambda_n^3 \hookrightarrow \Delta^3 \text{ rather than } I^3 \hookrightarrow \Delta^3.$$

Then we have $X \xrightarrow{\psi} NC$ defined as

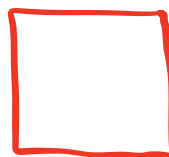
$$\begin{cases} X_0 \simeq (NC)_0 \\ X_n \simeq (NC)_n. \end{cases}$$

We prove that ψ_n is a bijection

by induction on $n \geq 2$.

$$\begin{array}{ccc} X_n & \xrightarrow[\sim]{(4)} & \text{sSet}(\Lambda_n^n, \mathbb{C}) \\ \downarrow & \equiv & \downarrow \textcircled{\star} \\ (\mathbb{N} \subset \mathbb{C})_n & \xrightarrow[\sim]{(1) \Rightarrow (4)} & \text{sSet}(\Lambda_n^n, \mathbb{N} \subset \mathbb{C}) \end{array}$$

But $\Lambda_n^n = \text{Sk}_{n-1}(\Lambda_n^n)$ is colimit of standard simplices of dimension $< n$, so $\textcircled{\star}$ is a bijection by induction.



Rmks

• We did not use the full strength of (4), only $\Lambda^1 \leftrightarrow \Delta^1$.

For later purposes it is (4) which is relevant.

• Moreover, we saw in the course of the proof that NC has the lifting property for $\Lambda^h \hookrightarrow \Delta^h$, $\underline{0} \leq h \leq \underline{n}$, as long as $n \geq 4$.

Prop 4

NC Kan complex



groupoid

\Leftrightarrow unique extension
prop for all horns.



proof: \Downarrow : Let $g: X \rightarrow Y$ be
a morphism in C . We look at

$$\Delta_0^2 \longrightarrow NC \quad g \nearrow Y$$

given by: $X \underset{id_X}{=} X$

By the Kan lifting property,
there exists $\sigma \in (NC)_2$

with
$$\begin{array}{ccc} g \nearrow Y & & \\ \sigma \searrow & & \\ X \underset{id_X}{=} X & & \end{array} \quad g := d^0(\sigma)$$

but then $\sigma \circ \rho = id$.

Similarly, using $\Lambda_2^2 \hookrightarrow \Delta^2$,
 $\exists g' : Y \rightarrow X$, $f \circ g' = \text{id}_Y$.

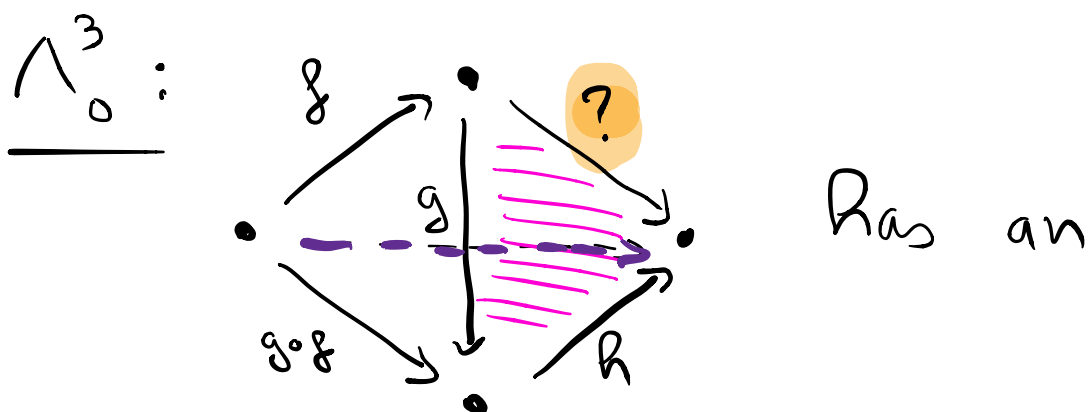
But then $g = g'$ and f is
 an isomorphism.

$$(g = g \circ (f \circ g') = (g \circ f) \circ g' = g') \\ \Rightarrow \subset \text{groupoid.}$$

\Uparrow : By Thm 34 and the
 remark above,

it is enough to prove the

lifting property for Λ_0^2 (Λ_2^2),
 Λ_0^3 (and Λ_3^3).



extension to $\Delta^3 \dashv NC$

iff $\text{?} = hg$.

By considering the other face, we see that

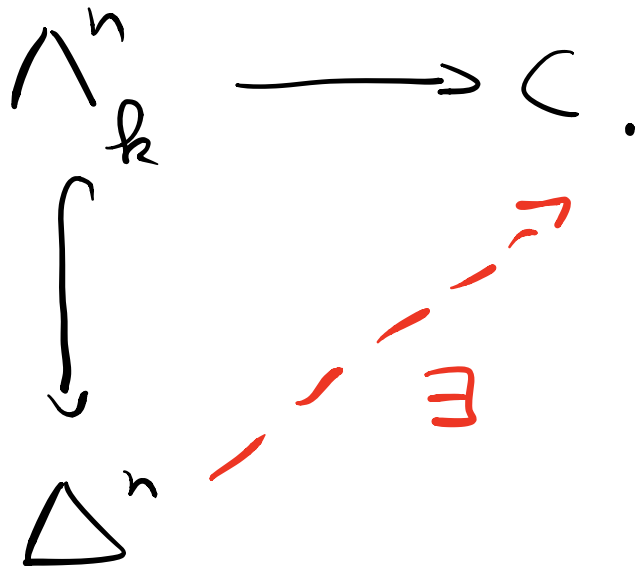
$$\text{?} \circ f = h(gf)$$

which is enough since f is iso. \square

5) ∞ -categories

def 5 An ∞ -category (or quasicategory) is a simplicial set $C \in \mathbf{sSet}$ satisfying the inner horn extension property:

$$\forall n \geq 2, \forall 0 < k < n,$$



- A **functor** $F: C. \rightarrow D.$ between ∞ -categories is simply a morphism of simplicial sets. This defines a (1-)category \mathbf{Cat}_{∞}^1 of ∞ -categories: $\mathbf{Cat}_{\infty}^1 \overset{\text{full}}{\longleftrightarrow} \mathbf{sSet}.$

- A **natural transformation** $C. \begin{matrix} \xrightarrow{F} \\ \alpha \parallel \\ \xrightarrow{G} \end{matrix} D.$ is a morphism $\alpha: C. \times \Delta^1 \rightarrow D.$ with $\alpha|_{C \times \{0\}} = F$ and $\alpha|_{C \times \{1\}} = G.$ □

Basic examples:

- By $\left\{ \begin{array}{l} \text{the def. of Kan complexes} \\ \text{Prop 33 and Thm 34} \end{array} \right.$, we have fully faithful functors:

$$\begin{array}{ccccc} & & \mathbf{Cat} & \xrightarrow{N} & \\ \mathbf{Grp} & \nearrow & & & \\ & & \mathbf{Kan} & \searrow & \\ & & & & \mathbf{Cat}_{\infty}^1 \longleftrightarrow \mathbf{sSet}. \end{array}$$

- We will see later that there are many examples which are not of these forms.

History

- This definition is due to Boardman-Vogt

(1973) in the context of homotopy theory (homotopy coherent algebraic structures, infinite loop spaces). They proved some basic results which we will review next.

- The idea of taking quasicategories as a model for $(\infty, 1)$ -categories is due to Joyal (late 90's) and he developed most of the results from in the first

half of this course. Then Lurie came and pushed the theory even further!

Terminology For $X \in \mathbf{sSet}$ (and in particular for ∞ -categories), we call

- **objects of X** . the elements of X_0

- **$(1-)$ morphisms of X** . $\text{---} X_1$

For $f \in X_1$, we say that the **source**

(resp. the **target**) of f is $d_1(f)$ (resp. $d_0(f)$)

and we write $f: d_1(f) \rightarrow d_0(f)$.

- For $x \in X_0$, we write $\text{id}_x = \Delta_0(x)$ and call it the **identity morphism** of x .

Slogan: An ∞ -category is like the nerve of a category, except that composition of chains of composable morphisms is only well-defined up to homotopy, and these homotopies are compatible in a precise way.

In particular, it should be possible to get a 1-category by identifying all those homotopies.

This amounts to giving a simple description of the fundamental category $\pi_1 X$ when X is an ∞ -category.

This was achieved by Boardman-Vogt.

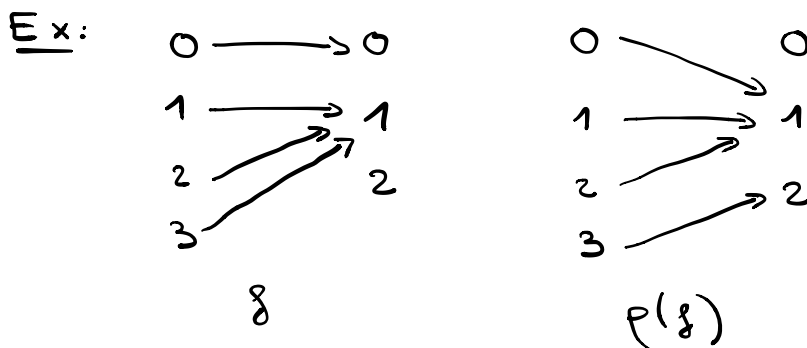
The first tool we want to have in any "category theory" is duality.

def 6 The **order-reversing** functor

$$p: \Delta \longrightarrow \Delta \text{ is defined as}$$

the identity on objects and, for $f: [m] \rightarrow [n]$,

$$p(f)(i) := n - f(m - i).$$



$p^*: \mathbf{sSet} \longrightarrow \mathbf{sSet}$ is the functor "precomposition

by p). For $X. \in \mathbf{sSet}$, the **opposite simplicial**

set is $X.^{op} := p^*(X.)$. □

Examples . $(\Delta^n)^{op} \cong \Delta^n$, $(I^n)^{op} \cong I^n$, $(\partial \Delta^n)^{op} \cong \partial \Delta^n$.

. $(\wedge_i^n)^{op} \cong \wedge_{n-i}^n$

. $X.$ ∞ -category (resp Kan)

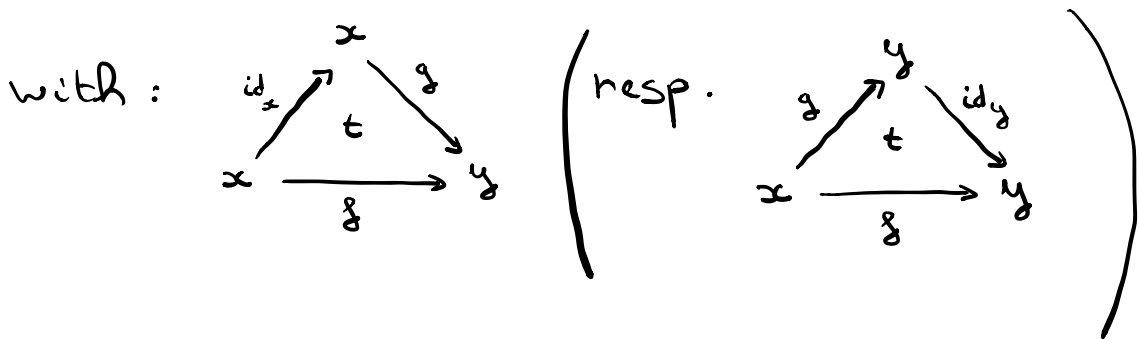
$\Leftrightarrow X.^{op}$ ∞ -category (resp Kan).

• $N(C^{op}) \cong N(C)$ for $C \in \text{Cat}$.

We now turn to the Boardman-Vogt result.

Def 7: Let $X_0 \in \text{sSet}$, $x, y \in X_0$.

We say that $f, g: x \rightarrow y$ are **left homotopic** (resp. **right homotopic**), written $f \sim_l g$ (resp. $f \sim_r g$) if $\exists t \in X_2$



Left and right homotopy are not necessarily equivalence relations; however we have

Lemma 8 Let C be an ∞ -category. Then

left and right homotopy coincide and is an equivalence relation.

Proof: Let $f, g, h: x \rightarrow y$ in $C(x, y)$.

We prove:

a) $f \underset{e}{\sim} f$.

b) $f \underset{e}{\sim} g$ and $g \underset{e}{\sim} h$ imply $f \underset{e}{\sim} h$.

c) $f \underset{e}{\sim} g$ implies $f \underset{r}{\sim} g$.

d) $f \underset{r}{\sim} g$ implies $g \underset{e}{\sim} f$.

a): $t := f_{001}$ works:

b), c), d): They are proven in the same way:

- construct from the given 2-simplices a map

$$\bigwedge_i^3 \rightarrow C, \text{ with } i = \begin{cases} 1, & \text{for b) \& c)} \\ 2, & \text{for d)} \end{cases}$$

- appeal to the inner horn extension property

to get $\Delta^3 \rightarrow C$.

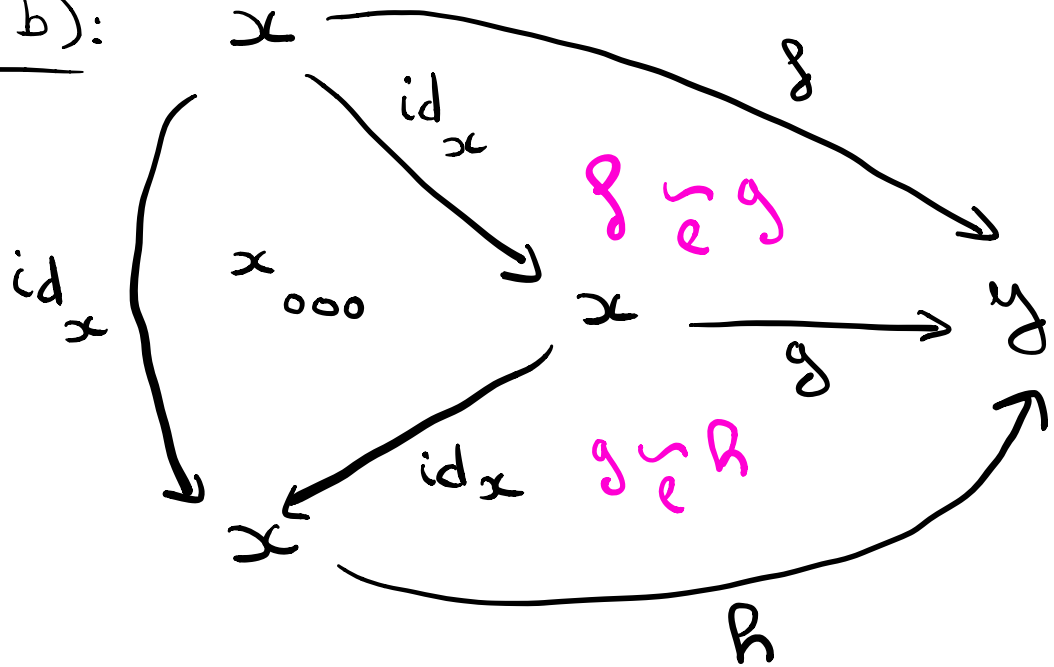
- restrict to the "new" 2-simplex.

It is easier in this case to display

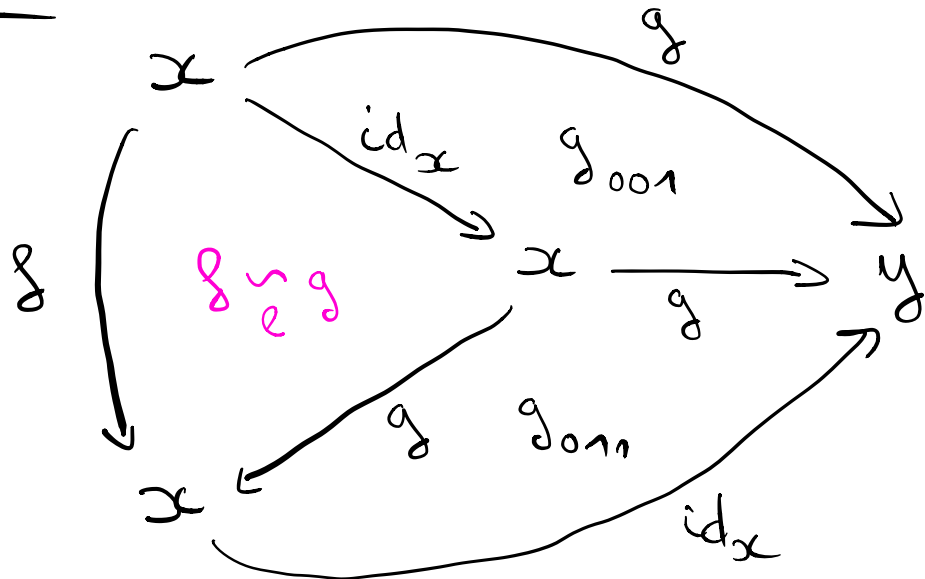
simplices so that the missing face is the back

of the picture:

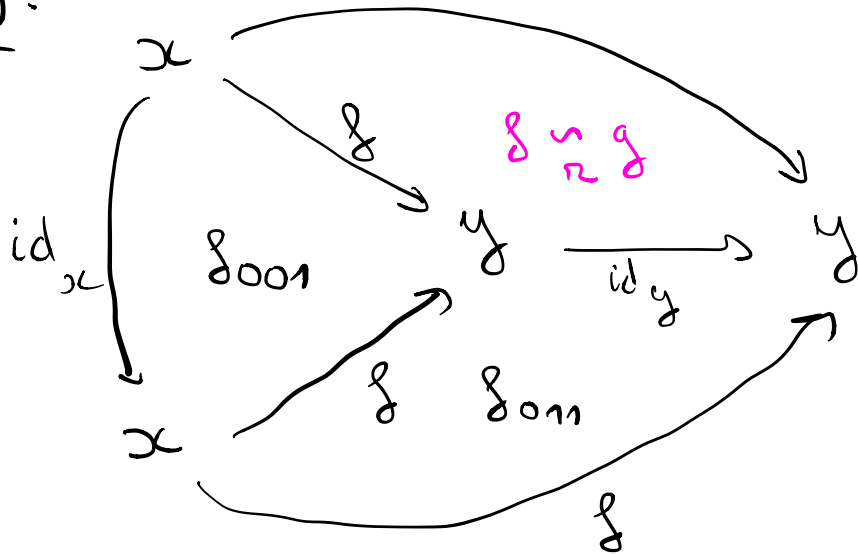
For b):



For c):



For d):



Finally:

- c), d) $\Rightarrow \simeq_e$ is symmetric.
- + a), b) $\Rightarrow \simeq_e$ equivalence relation
- + c), d) $\Rightarrow \simeq_e = \simeq_r$

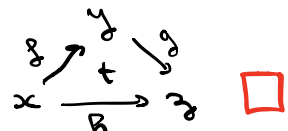


In this case we write $f \simeq_r g$ for $f \simeq_e g$.
and we write $[f]$ for the **homotopy class** of f in $C(x, y)$.

def 9: Let $C \in \text{Cat}_\infty$, $f: x \rightarrow y, g: y \rightarrow z$

and $h: x \rightarrow z$. We say that h is a **composition**

of f and g if there is $t \in C_2$ with

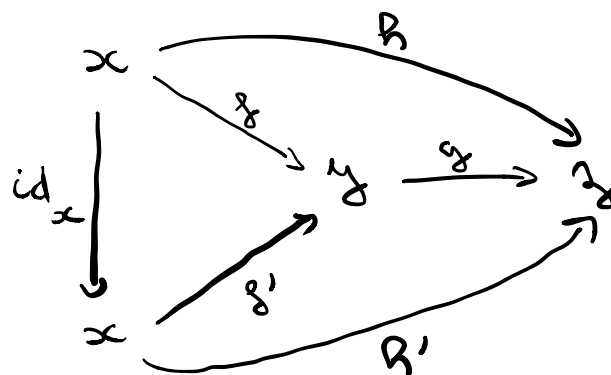


Prop 10: In an ∞ -category, compositions exist; their homotopy class is well-defined and depends only on the homotopy classes of the morphisms being composed.

proof: The existence is simply the extension property for $\Lambda_n^2 \hookrightarrow \Delta^2$.

Let $\begin{cases} f \simeq f': x \rightarrow y \\ g \simeq g': x' \rightarrow y' \end{cases}$ and let $\begin{cases} R & \text{be a composition of } f \text{ and } g \\ R' & \text{be a composition of } f' \text{ and } g'. \end{cases}$

We must prove that $R \simeq R'$. It is enough to treat separately the cases $f = f'$ and $g = g'$. By working in C^{op} , we reduce to $f = f'$. As before we construct a horn Λ_2^3 , extends and restrict:



Prop 11: The resulting composition on homotopy classes of morphisms is associative and unital.

proof: same method as Prop 39: left as exercise \square

def 12 Let C be an ∞ -category. Its homotopy category hC has:

$$\begin{cases} \text{Ob}(hC) = C_0 \\ hC(x, y) = C(x, y) / \simeq \end{cases} \quad \begin{array}{l} \text{By Prop 39-40, this} \\ \text{is indeed a 1-category.} \end{array}$$

Prop 13: The construction of the homotopy category defines a functor $R: \text{Cat}_\infty^1 \rightarrow \text{Cat}$ which is a left adjoint to $N: \text{Cat} \rightarrow \text{Cat}_\infty^1$.

(\Leftarrow) $\tau C \simeq hC$ naturally in C

proof: The functoriality follows from the fact that for $F: C \rightarrow D$, $f \simeq f'$ in $C \Rightarrow F(f) \simeq F(f')$ which is clear.

Let C be an ∞ -category. We construct a natural equivalence (in fact isomorphism) of categories $\tau C \xrightarrow[\varphi]{\simeq} hC$.

Because of Prop 3, to define φ it is

enough to define
$$\begin{cases} C_0 \rightarrow \text{Ob } \mathcal{R}C \\ C_1 \rightarrow \text{Mor } \mathcal{R}C \end{cases}$$

Satisfying some relations given by identities and C_2 .

We put $C_0 = \text{Ob } \mathcal{R}C$

$C_1 \rightarrow \text{Mor } \mathcal{R}C, f \mapsto [f]$

and the relations are satisfied by def 12.

- By construction, φ is $\begin{cases} \text{bijective on objects} \\ \text{surjective on morphisms.} \end{cases}$

It remains to show φ is faithful.

Because of the lifting property for $\Lambda^2 \hookrightarrow \Delta^2$,

any morphism in πC can be written as \overline{f}

for $f \in C_1$. Now suppose that $f, f': x \rightarrow y$

satisfy $[f] = [f']$. By definition, there is

an homotopy
$$\begin{array}{ccc} & x & \\ & \parallel & \searrow f' \\ x & \xrightarrow{f} & y \end{array}$$
, but this also

implies $\text{id}_x \circ \overline{f'} = \overline{f}$ in πC and we are done. \square

Let us see some basic ways to construct new ∞ -categories.

Prop 14: 1) Arbitrary products and coproducts of ∞ -categories (in $sSet$) are ∞ -categories.

2) Filtered colimits of ∞ -categories are ∞ -categories.

proof: 1)

Let $\{X_\alpha\}_{\alpha \in J}$ be a family of ∞ -categories.

. Let $0 \leq h \leq n$. We have

$$sSet\left(\Delta^n, \prod_{\alpha} X_{\alpha}\right) \longrightarrow sSet\left(\Lambda_{\mathbb{R}}^n, \prod_{\alpha} X_{\alpha}\right)$$

IS IS

$$\prod_{\alpha} sSet(\Delta^n, X_{\alpha}) \longrightarrow \prod_{\alpha} sSet(\Lambda_{\mathbb{R}}^n, X_{\alpha})$$

↑
product of
surjections is
a surjection

$\Rightarrow \prod_{\alpha} X_{\alpha}$ is an ∞ -category.

- For the coproduct, we need to compute $s\text{Set}(\Lambda_R^n, \coprod X_\alpha)$ and $s\text{Set}(\Delta^n, \coprod X_\alpha)$

By Yoneda, $s\text{Set}(\Delta^n, \coprod X_\alpha) = (\coprod X_\alpha)([n])$

colimits
are objectwise

$$= \coprod X_\alpha([n])$$

$$= \coprod s\text{Set}(\Delta^n, X_\alpha)$$

So it suffices to show that the natural

map

$$\coprod s\text{Set}(\Lambda_R^n, X_\alpha) \longrightarrow s\text{Set}(\Lambda_R^n, \coprod X_\alpha)$$

is a bijection. For this one can use

$$\Lambda_R^n = \coprod_{\substack{\Delta^{[n]-\{i,j\}} \\ i \neq j}} \Delta^{[n]-i} \quad \text{and the Yoneda}$$

trick above; the key point is that the various $\Delta^{[n]-i}$ must be sent to the same X_α because they are connected via

the $(n-2)$ -faces $\Delta^{[n]-\{i,j\}}$. Or in other words, one can show that

$\pi_0(\coprod X_\alpha) \cong \coprod X_\alpha$ while $\pi_0(\Lambda_R^n)$ has one element (Λ_R^n is connected).

2) Let J be a filtered category and $X: J \rightarrow \mathbf{sSet}$ be a diagram so that each $X(\alpha)$ is an ∞ -category. Once again we have $(\operatorname{colim}_J X)_n = \operatorname{colim}_J X(n)$ because colimits are computed objectwise.

We want to show that the canonical map

$$\operatorname{colim}_{\alpha \in J} \mathbf{sSet}(\Lambda_R^n, X(\alpha)) \rightarrow \mathbf{sSet}(\Lambda_R^n, \operatorname{colim}_J X)$$

is a bijection. We will show this holds for Λ_R^n replaced by any $Y \in \mathbf{sSet}$ with finitely many non-degenerate simplices.

Let $\mathcal{C} = \left\{ \gamma \in \text{sSet} \mid \begin{array}{l} \text{sSet}(\gamma, -) \text{ commutes} \\ \text{with filtered colimits} \end{array} \right\}$.

As remarked above, $\Delta^n \in \mathcal{C}$ for all $n \in \mathbb{N}$.

Let's show that \mathcal{C} is closed under finite colimits. Let K be a finite category and

$\gamma : K \rightarrow \mathcal{C}$ be a diagram.

We have

$$\begin{array}{ccc}
 \text{Colim}_{\alpha \in J} \text{sSet} \left(\text{Colim}_K \gamma, X(\alpha) \right) & \longrightarrow & \text{sSet} \left(\text{Colim}_K \gamma, \text{Colim}_J X \right) \\
 \downarrow \text{IS colim prop.} & & \downarrow \text{IS colim prop.} \\
 \text{Colim}_{\alpha \in J} \text{Lim}_{\beta \in K} \text{sSet}(\gamma(\beta), X(\alpha)) & & \text{Lim}_{\beta \in K} \text{sSet}(\gamma(\beta), \text{Colim}_J X) \\
 \downarrow \text{Filtered colimits commutes with finite limits in Set.} & \searrow & \downarrow \text{IS } X(\alpha) \in \mathcal{C} \\
 & & \text{Lim}_{\beta \in K} \text{Colim}_{\alpha \in J} \text{sSet}(\gamma(\beta), X(\alpha))
 \end{array}$$

$\Rightarrow \text{Colim}_K \gamma \in \mathcal{C}$.

• By the skeletal filtration,

$$\left\{ \begin{array}{l} \text{s. sets with } < \infty \\ \text{non-deg. simplices} \end{array} \right\} = \left\{ \begin{array}{l} \text{finite colimits of} \\ \text{standard simplices} \end{array} \right\}$$

and we are done.

